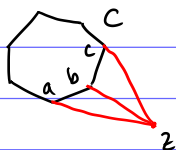


Q6.  $G$  nonbipartite,  $\Delta$ -free simple graph  $n$  vertices  $\delta(G) = k$

$\ell$  is min length of an odd cycle

a)



Suppose  $\exists z$  with 3 neighbours in  $C$ .

Then it forms three cycles

$$zab \quad zbc \quad zca$$

Assume  $a, b, c$  appear in order along the cycle. Let  $d(a, b)$  represent the # of edges between  $a$  and  $b$  in the cyclic order

$$d(a, b) + d(b, c) + d(c, a) = \ell$$

At least one them must be odd, since  $\ell$  is odd. WLOG let  $d(a, b)$  be odd. Then  $zab$  is an

So we must have  $2 + d(a, b) \geq \ell = d(a, b) + d(b, c) + d(c, a)$

This is only possible if  $d(a, b) = d(b, c) = 1 \Rightarrow 2$  triangles. This is a contradiction.

Note: The oddness of  $\ell$  is not that critical here. So I don't think it's required. So let  $\ell$  be a minimal cycle of any length.

Then the cycles have length

$$2 + d(a, b)$$

$$2 + d(b, c)$$

$$2 + d(c, a)$$

They all have length at least  $\ell$  (by definition of  $\ell$ )

$$\Rightarrow 2 + d(a, b) + 2 + d(b, c) + 2 + d(c, a) \geq 3\ell$$

But this gives us

$$6 + \ell \geq 3\ell \Rightarrow \ell \leq 3.$$

But our graph is triangle free, so this is a contradiction.

(We cannot have  $\ell = 2$  or  $\ell = 1$  since the graph is simple)

b)

# of edges between  $V(C)$  and  $V(G) - V(C)$ .

$$= \sum_{v \in V(C)} d(v) - \overset{\text{\# of edges within } V(C)}{\underbrace{2}} = |V(C)| (k-2) = e(k-2)$$

Since each  $v \in V(G) - V(C)$  can be connected to at most 2  $V(C)$  vertices

we also have the upper bound

$$\leq 2 |V(G) - V(C)| = 2(n-e)$$

$$\Rightarrow e(k-2) \leq 2(n-e) \Rightarrow \frac{ke}{2} \leq n$$

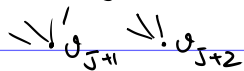
Q7 a) All vertices have even degree.

Prob 1.2.27 : Every even graph decomposes into cycles.

$$G = C_1 \sqcup C_2 \dots \sqcup C_k$$

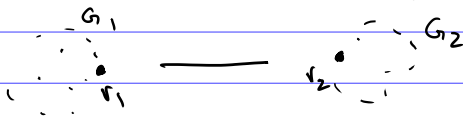
Take any edge, it must belong to one of the cycles  $\Rightarrow$  it cannot be a cut edge

b)  $v_1 \dots v_j \quad j-2$



It's obvious.

c) Take 2 copies of a graph



$G_1$  has the following props

1.  $d(r_1) = 2k$

2.  $d(v) = 2k+1 \quad v \in G - r_1$

Connect  $r_1$  and  $r_2$ . Now  $d(r_1) = d(r_2) = 2k+1$ .  $r_1, r_2$  is a cut edge.

How to construct  $G_1$ ?

Form a  $2k$  regular graph with  $2k+2$  vertices.  $v_1, \dots, v_{2k+1}, v_{2k+2}$  where the last two vertices are not connected to each other — this by part (b) and an obvious induction.

and

Create  $r_1$ , connect it to  $2k$  vertices, and connect the 2 remaining vertices  $v_{2k+1}, v_{2k+2}$ .

$$\deg(r_1) = 2k, \quad d(v_i) = 2k+1$$

■

